Internet Protocols

Supporting Protocols and Framing

Updated: 9/30/14

Supporting Protocols

ARP / RARP
BOOTP
ICMP
DHCP
NAT

IP Supporting Protocols

- IP protocol only deals with the data transfer (best-effort)
 - Possible Errors that can happen and not detected by IP: Data lost, duplication, out-oforder
 - However there are some error checking mechanisms:
 - CRC, TTL

IP Supporting Protocols

We focus on the following Protocols: ICMP, ARP, RARP, BOOTP, DHCP





- Used to resolve network layer addresses into link layer addresses
- Exploits broadcast property of a LAN
- Each host on LAN maintains a table of IP subnetwork addresses
- If the address can not be found ARP broadcasts a request
 - Shouting: Who knows about this IP address?
- Other hosts listen and reply
 - The reply includes IP address and MAC (unicast)
 - Any interested host can learn about the new information



- Assume **1** is sending a message to **2** (192.31.65.5)
 - What is the MAC address for 192.31.65.5? Use ARP broadcast!
 - Host 2 responds to Host 1: it is E2
 - Host 1 maps IP and MAC;
 - Encapsulate the IP message in the Ethernet frame and sends it
 - Cashing can enhance ARP operation (Node 1 can cash the result)



- Assume **1** is sending a message to **4** (<u>rose@ee.sonoma.edu</u>)
 - <u>ee.sonoma.edu</u> is the destination

ARP Example

- Host 1 sends a message to Domain Name System (DNS): what is the IP address for <u>ee.sonoma.edu</u>? → 192.31.63.8
- What is the MAC address for 192.31.63.8? ARP cannot pass through the router!
- Two choices:
 - 1. Reconfigure routers to respond to ARP (Proxy ARP)
 - The ARP Proxy is aware of the location of the destination
 - Proxy offers its own MAC address
 - Thus, it acts on behalf of the node: "send it to me, and I'll get it to where it needs to go."
 - In this example the Proxy can be E4
 - 2. Send the message to the LAN router
 - Note that ARP is limited to a single network
 - In the example above, the address binding or resolution is done between Node 1 and E3; then between E3 and E4; then E4 and node 4 (via broadcast).
 - Node 4 will send back its MAC to node 1 (not found in ARP cache)
 - Each router looks at the IP address and passes it to the next node using the routing table

ARP <u>Request</u> Content - Broadcast

<u>F</u> ilter:			▼ <u>E</u> xpression (lear <u>A</u> pply	
No. +	Time	Source	Destination	Protocol	Info
	1 00:19:20.15713	O AmbitMic_a9:3d:68	Broadcast	ARP	Who has 192.168.1.1? Tell 192.168.1.105
	2 00:19:20.15814	8 LinksysG_da:af:73	AmbitMic_a9:3d:68	ARP	192.168.1.1 is at 00:06:25:da:af:73
	3 00:19:20.15815	8 AmbitMic_a9:3d:68	LinksysG_da:af:73	0x0800	IP
	4 00:19:23.11998	0 AmbitMic_a9:3d:68	LinksysG_da:af:73	0x0800	IP
	5 00:19:29.12861	.8 AmbitMic_a9:3d:68	LinksysG_da:af:73	0x0800	IP
	6 00:19:33.70010	4 Telebit_73:8d:ce	Broadcast	ARP	who has 192.168.1.117? Tell 192.168.1.104
	7 00:19:37.60155	3 AmbitMic_a9:3d:68	LinksysG_da:af:73	0x0800	IP
	8 00:19:37.62303	2 LinksysG_da:af:73	AmbitMic_a9:3d:68	0x0800	IP
	9 00:19:37.62305	7 AmbitMic_a9:3d:68	LinksysG_da:af:73	0x0800	IP
	10 00:19:37.62359	08 AmbitMic_a9:3d:68	LinksysG_da:af:73	0x0800	IP
	11 00:19:37.65189	6 LinksysG_da:af:/3	AmbitMic_a9:3d:68	0x0800	IP
	ernet II, Src: Am Destination: Broad Address: Broadca 1	<pre>bitMic_a9:3d:68 (00:c lcast (ff:ff:ff:ff:ff: ast (ff:ff:ff:ff:ff:ff: = IG bi = LG bi = LG bi</pre>	0:59:a9:3d:68), Dst: B ff) Destination t: Group address (mult t: Locally administere	noadcast (on add icast/broa d address	(ff:ff:ff:ff:ff:ff) ress adcast) (this is NOT the factory default)
	Address: AmbitMi	ic_a9:3d:68 (00:d0:59: IG b1 LG b1	a9:3d:68) Sourc t: Individual address t: Globally unique add	e addr (unicast) Iress (fact	ess ory default)
Т	ype: ARP (0x0806))			
Add	ress Resolution P	Protocol (request)			
H	ardware type: Eth	nernet (0x0001)			
F	protocol type: IP	(0x0800)			
H	ardware size: 6				
	rotocol size: 4				
	neodo: noquoct ((>0001)			
0000	<i><i><i><i>i</i> i i i i i i i i</i> </i></i>	E 00 d0 50 a0 3d 68 a	08.06.00.01	v _h	
0010	08 00 06 04 00 0	1 00 d0 59 a9 50 68	0 28 01 69	Y =h 1	
OUTO .					

ARP Request Content – Contains IP Address

Timer:			• <u>Expression</u>	lear <u>A</u> pply	
No	Time	Source	Destination	Protocol	Info
	1 00:19:20.157130	AmbitMic_a9:3d:68	Broadcast	ARP	who has 192.168.1.1? Tell 192.168.1.105
	2 00:19:20.158148	LinksysG_da:af:73	AmbitMic_a9:3d:68	ARP	192.168.1.1 is at 00:06:25:da:af:73
	3 00:19:20.158158	AmbitMic_a9:3d:68	LinksysG_da:af:73	0x0800	IP
1	4 00:19:23.119980	AmbitMic_a9:3d:68	LinksysG_da:af:73	0x0800	IP
	5 00:19:29.128618	AmbitMic_a9:3d:68	LinksysG_da:af:73	0x0800	IP
	7 00:10:27 601552	Teleoit_/3:80:ce	Broducast	ARP 0x0800	who has 192.168.1.11/? Tell 192.168.1.104
1	8 00:10:37 623032	LinksysC dataf:73	AmbitMic 20:3d:68	0x0800	19
	9 00:19:37 623057	AmbitMic a9:3d:68	Linksysc da:af:73	0x0800	TP
	10 00:19:37.623598	AmbitMic a9:3d:68	LinksysG da:af:73	0x0800	IP
	11 00:19:37.651896	LinksysG da:af:73	AmbitMic_a9:3d:68	0x0800	IP
1	12 00:19:37.656065	LinksvsG da:af:73	AmbitMic a9:3d:68	0x0800	TP
⊟ Ad	Trailer: 0000000000 dress Resolution Pro	00000000000000000000000000000000000000	00000		
	Haroware type: Ethe	rhet (0x0001)			
	Protocol type: IP (0x0800)			
	Haroware size: 6		ADD massage	conto	n the ID address of the conden
۱.	Protocol size: 4		AKP message	conta	IT THE LP ADDRESS OF THE SENDER
1	FIOLOCOT 312E. 4				in the fr dutiess of the sender
(Opcode: request (Ox	0001)			in the fr douress of the sender
	opcode: request (0x Sender MAC address:	0001) Telebit_72:8d:ce (00	1.80.ad.72.8d:ce)		in the fr dataress of the sender
	opcode: request (0x) Sender MAC address: Sender IP address: 3	0001) Telebit_72:8d:ce (00 192.168.1.104 (192.16	580-ad-72-8d:ce)		in the fr data cost of the sender
Ę	opcode: request (0x) Sender MAC address: Sender IP address: Target MAC address:	0001) Telebit_72:8d:ce_(00 192.168.1.104_(192.16 00:00:00_00:00:00 (0	0:00:0d:72:8d:ce) (8.1.104) (0:00:00:00:00:00)		
Ę	opcode: request (0x) Sender MAC address: Sender IP address: Target MAC address: Target IP address:	0001) Telebit 72:8d:ce (00 192.168.1.104 (192.16 00:00:00_00:00:00 (0 192.168.1.117 (192.16	0:00:0d:72:8d:ce) 58.1.104) 0:00:00:00:00:00:00) 58.1.117)		
	opcode: request (0x) Sender MAC address: Sender IP address: Target MAC address: Target IP address:	0001) Telebit 72:8d:ce (00 192.168.1.104 (192.16 00:00:00_00:00:00 (0 192.168.1.117 (192.16 00.80 ad 73 8d ce 0	580-3d-72-8d:ce) 58.1.104) 00:00:00:00:00:00) 58.1.117)		
0000	opcode: request (0x) Sender MAC address: Sender IP address: Target MAC address: Target IP address: ff ff ff ff ff ff ff 08 00 06 04 00 01	0001) Telebit_72:8d:ce_00 192.168.1.104 (192.16 00:00:00_00:00:00 (0 192.168.1.117 (192.16 00 80 ad 73 8d ce 0 00 80 ad 73 8d ce 0	80.3d.72.8d:ce) 58.1.104) 00:00:00:00:00:00) 58.1.117) 8 06 00 01 0 28 01 58	.s	
0000	opcode: request (0x) Sender IP address: Target MAC address: Target IP address: Target IP address: ff ff ff ff ff ff 08 00 06 04 00 01 00 00 00 00 00 00	0001) Telebit_72:8d:ce_(00 192.168.1.104 (192.16 00:00:00_00:00:00 (0 192.168.1.117 (192.16 00 80 ad 73 8d ce 0 00 80 ad 73 8d ce c co a8 01 75 00 000	80.3d.72.8d:ce) 58.1.104) 50:00:00:00:00:00) 58.1.117) 8 06 00 01 0 a8 01 68	.sh .sh	

ARP Message Format



ARP Message Format

HARDWARE ADDRESS TYPE

- 16-bit field that specifies the type of hardware address being used
- the value is 1 for Ethernet
- PROTOCOL ADDRESS TYPE
 - 16-bit field that specifies the type of protocol address being used
 - the value is 0x0800 for IPv4
- HADDR LEN
 - 8-bit integer that specifies the size of a hardware address in bytes
- PADDR LEN
 - 8-bit integer that specifies the size of a protocol address in bytes
- OPERATION
 - 16-bit field that specifies whether the message
 - request (the field contains 1) or
 - response (the field contains 2)

0		8	16	24	31	
	HARDWARE ADDRESS TYPE			ROTOCOL ADDRESS TYPE		
	HADDR LEN	PADDR LEN		OPERATION		
	SENDER HADDR (first 4 octets)					
	SENDER HADD	R (last 2 octets)	SE	NDER PADDR (first 2 octets)		
	SENDER PADD	R (last 2 octets)	ТА	RGET HADDR (first 2 octets)		
		TARGET HADD	R (last	4 octets)		
		TARGET PAD	DR (all 4	octets)		



ARP Message Format

0		8	16	24	31
	HARDWARE ADDRESS TYPE			ROTOCOL ADDRESS TYPE	
	HADDR LEN	PADDR LEN		OPERATION	
		SENDER HADD	R (first	4 octets)	
	SENDER HADD	R (last 2 octets)	SE	NDER PADDR (first 2 octets)	
	SENDER PADD	R (last 2 octets)	ТА	RGET HADDR (first 2 octets)	
		TARGET HADD	R (last	4 octets)	
		TARGET PAD	DR (all	4 octets)	

- SENDER HADDR
 - HADDR LEN bytes for the sender's hardware address
- SENDER PADDR
 - PADDR LEN bytes for the sender's protocol address
- TARGET HADDR
 - HADDR LEN bytes for the target's hardware address
- TARGET PADDR
 - PADDR LEN bytes for the target's protocol address

```
Address Resolution Protocol (request)
Hardware type: Ethernet (0x0001)
Protocol type: IP (0x0800)
Hardware size: 6
Protocol size: 4
Opcode: request (0x0001)
Sender Mac address: Telebit 72:8d:ce (00:80:ad:72:8d:ce)
Sender IP address: 192.168.1.104 (192.168.1.104)
Target MAC address: 00:00:00_00:00:00 (00:00:00:00:00)
Target IP address: 192.168.1.117 (192.168.1.117)
```

Notes

ARP is encapsulated in Ethernet frame

- In this case Ethernet type will be ARP
- Sending ARP for each message is not efficient
 - Thus, cache is used (create a small local table)
 - The cache is checked before broadcasting the request



Internet Control Message Protocol (ICMP)



- ICMP error messages are used by routers and hosts to tell a device that sent a datagram about problems encountered in delivering it
 - It is used to test the network
 - ICMP messages are encapsulated in the IP packet
 - ICMP has many message types
 - Two basic categories: Report Error or Obtain Information

Number	Туре	Purpose
0	Echo Reply	Used by the ping program
3	Dest. Unreachable	Datagram could not be delivered
5	Redirect	Host must change a route
8	Echo	Used by the ping program
11	Time Exceeded	TTL expired or fragments timed out
12	Parameter Problem	IP header is incorrect
30	Traceroute	Used by the traceroute program

IP ICMP Hdr ICMP Payload

MIME

HTTP SMTP TELNET

TCP

SNMP

UDP

ICMP IGMP OSPF RSVP

BGP

FTP

Code field is used for subtypes

ICMP Protocol Details



- **Type:** 8 bits
- Code: 8 bits
- ICMP Header Checksum: 16 bits
 - The 16-bit one's complement of the one's complement sum of the ICMP message $(6 \rightarrow -6)$
- Data: Variable length
 - Contains the data specific to the message type indicated by the Type and Code fields.

Encapsulated ICMP – Type 8

🕲 t.pcap - Ethereal				
Elle Edit View Go Capture Anal	yze Statistics Help			
	🖌 🗙 🖓 📇		⇒ 🕫 🛠 🖗 🖨	
Elter: ip.addr == 192.168.2.157 &8	icmp	•	Expression Glear Apply	
No. + Source	Destination	Protocol Info		Deita Time
6086 192.168.2.157	66.249.93.104	ICMP Ech	o (ping) request	3211.32393
6088 66.249.93.104	192.168.2.157	ICMP Ech	o (ping) reply	0.076788
6089 192.168.2.157	66.249.93.104	ICMP Ech	o (ping) request	0.925777
6090 66.249.93.104	192.168.2.157	ICMP Echi	o (ping) reply	0.076979
6094 192.168.2.157	66.249.93.104	ICMP Ech	o (ping) request	0.924109
6096 66.249.93.104	192.168.2.157	ICMP Ech	o (ping) reply	0.075390
6098 192.168.2.157	66.249.93.104	ICMP Ech	o (ping) request	0.925826
6099 66.249.93.104	103 168 3 157	TOMO Eshi	vines canly	0.075460



Encapsulated ICMP – Type 8

```
Internet Protocol, Src: 192.168.2.157 (192.168.2.157), Dst: 66.249.93.104 (66.249.93.104)
   Version: 4
   Header length: 20 bytes
   Differentiated Services Field: 0x00 (DSCP 0x00: Default; ECN: 0x00)
       0000 00.. = Differentiated Services Codepoint: Default (0x00)
       .... ..0. = ECN-Capable Transport (ECT): 0
       \dots 0 = ECN-CE: 0
   Total Length: 60
   Identification: 0x0fd3 (4051)
                                          Internet Control Message Protocol
   Flags: 0x00
       0... = Reserved bit: Not set
                                               Type: 8 (Echo (ping) request)
       .0.. = Don't fragment: Not set
                                               Code: 0
                                           8 Bytes
       Checksum: 0x475c [correct]
   Fragment offset: 0
   Time to live: 128
                                               Identifier: 0x0300
   Protocol: ICMP (0x01)
                                               Sequence number: 0x0300
   Header checksum: 0xc747 [correct]
       Good: True
                                               Data (32 bytes)
       Bad : False
   Source: 192.168.2.157 (192.168.2.157)
   Destination: 66.249.93.104 (66.249.93.104)
```

Type. 8 bits. Set to 8. Code. 8 bits. Cleared to 0. ICMP Header Checksum. 16 bits. Identifier. 16 bits. This field is used to help match echo requests to the associated reply. It may be cleared to zero. Sequence number. 16 bits. This field is used to help match echo requests to the associated reply. It may be cleared to zero. Data. Variable length. Implementation specific data (depends on the type)

Number	Туре	Purpose
0	Echo Reply	Used by the ping program
3	Dest. Unreachable	Datagram could not be delivered
5	Redirect	Host must change a route
8	Echo	Used by the ping program
11	Time Exceeded	TTL expired or fragments timed out
12	Parameter Problem	IP header is incorrect
30	Traceroute	Used by the traceroute program

ICMP

ICMP messages do not have special priority

They are forwarded like any other datagram, with one minor exception:

If an ICMP error message causes an error no error message is sent

The reason should be clear:

the designers wanted to avoid the Internet becoming congested carrying error messages about error messages

ICMP & Traceroute

- Traceroute is a program that shows you route taken by packets through a network
- sends a UDP packet to the destination taking advantage of ICMP's messages



Protocol Software and Configuration



Protocol Software and Configuration

- Once a host or router has been powered on, OS is started and the protocol software is initialized How does the protocol software in a host or router begin operation?
- For a router, the configuration manager must specify initial values for items such as
 - the IP address for each network connection
 - the protocol software to run
 - initial values for a forwarding table
 - the configuration is saved, and a router loads the values during startup

Host configuration usually uses a two-step process, known as bootstrapping

- A protocol was invented to allow a host to obtain multiple parameters with a single request, known as the Bootstrap Protocol (BOOTP)
 - Examples of such parameters: IP address; MASK; Local DNS
- Currently, DHCP is used to take care of most configuration needed

RARP and BOOTP

- Reverse ARP translates the Ethernet address to IP address
 - A diskless machine when it is booting can ask: My MAC is 12.03.23.43.23.23; what is my IP?
- RARP broadcasts the question (destination address is <u>all one</u>)
 - Not passed through the router!
- Major issue: Each LAN needs a <u>RARP server</u>!
- Bootstrap protocol uses UDP and forwards over routers
 - BOOTP is usually used during the bootstrap process when a computer is starting up
 - Mapping must be done manually in each router!

Dynamic Host Configuration Protocol

- DHCP allows a computer to join a new network and obtain an IP address automatically
 - The concept has been termed plug-and-play networking
- Replaces BOOTP and RARP
 - Extension of BOOTP data format
- DHCP uses UDP
 - UDP port 67 for sending data to the server
 - UDP port 68 for data to the client
- DHCP communications are connectionless in nature

Dynamic Host Configuration Protocol

DHCP has four basic phases:

- IP discovery, IP lease offer, IP request, and IP lease acknowledgement
- First DHCP server must be discovered
 - The client broadcasts messages on the physical subnet to discover available DHCP servers
- IP Lease Offer
 - When a DHCP server receives an IP lease request from a client, it reserves an IP address for the client and extends an IP lease offer by sending a DHCP OFFER message to the client

ſ	No	Len	Time	Source	Destination	Protocol	Info			
l	1	314	0.000000	0.0.0.0	255.255.255.255	DHCP	DHCP Discover	-	Transaction	ID
ſ	2	342	0.000295	192.168.0.1	192.168.0.10	DHCP	DHCP Offer	-	Transaction	ID (
I	3	314	0.070031	0.0.0.0	255.255.255.255	DHCP	DHCP Request	-	Transaction	ID (
I	4	342	0.070345	192.168.0.1	192.168.0.10	DHCP	DHCP ACK	-	Transaction	ID (

Dynamic Host Configuration Protocol

- A client can receive multiple offers from difference servers
 - Thus, it must request an IP address
- DHCP sends a Request packet to the DHCP server and receives a DHCP Reply
 - What is the IP address for this MAC?
 - It can also request its previous IP address!
- Even when an IP address is assigned, how long is it good for?
 - Before the IP address is removed find another IP address....called Leasing
- When the DHCP server receives the Request from the client, the configuration process enters its final phase
 - a DHCPACK (ACK) packet is sent to the client

No	Len	Time	Source	Destination	Protocol	Info
1	314	0.000000	0.0.0.0	255.255.255.255	DHCP	DHCP Discover - Transaction ID
2	342	0.000295	192.168.0.1	192.168.0.10	DHCP	DHCP Offer - Transaction ID
3	314	0.070031	0.0.0.0	255.255.255.255	DHCP	DHCP Request - Transaction ID
4	342	0.070345	192.168.0.1	192.168.0.10	DHCP	DHCP ACK - Transaction ID

DHCP

- DHCP includes several important details that optimize performance, such as the following:
- Recovery from loss or duplication
 - DHCP is designed to insure that missing or duplicate packets do not result in misconfiguration
 - If no response is received
 - a host <u>retransmits</u> its request (remember DHCP uses UDP!)
 - If a duplicate response arrives
 - a host ignores the extra copy
- Caching of a server address
 - once a host finds a DHCP server
 - the host caches the server's address
- Avoidance of synchronized flooding
 - DCHP takes steps to prevent synchronized requests
 - Synchronization can occur when all computers boot up at the same time!



DHCP Phases

dhcp.pcap - Wireshark	
<u>File Edit View Go Capture Analyze Statistics Telephony</u> <u>T</u> ools <u>H</u> elp	
▋₩₩₩₩₩₽₽₩₩₩₽₽₩₩₩₽₽₩₩₽₽	-C
F <u>i</u> lter: ▼ E <u>x</u> pression Clea <u>r</u> App <u>ly</u>	
802.11 Channel: Channel Offset FCS Filter: Decryption Mode: Wireshark	
No Len Time Source Destination Protocol Info	
1 314 0.000000 0.0.0.0 255.255.255 DHCP DHCP Discover - Transaction ID	
2 342 0.000295 192.168.0.1 192.168.0.10 DHCP DHCP OTTER - Transaction ID 3 314 0.070031 0.0.0 255.255.255.255 DHCP DHCP Request - Transaction ID	
4 342 0.070345 192.168.0.1 192.168.0.10 DHCP DHCP ACK - Transaction ID	
0 8 16	24 31
OP HTYPE	HLEN HOPS
TRANSACTION ID	ENTIFIER
BOOTSTRAP PROTOCOT SECONDS ELAPSED	FLAGS
Message type: Boot Request (1) CLIENT IP ADD	DRESS
Hardware address length: 6	RESS
Hons: 0	DRESS
Transaction TD: 0x00003d1d	DRESS
Seconds elapsed: 0	RESS (16 OCTETS)
Bootp flags: 0x0000 (Unicast) Bootp flags: 0x000 (Unicast) Bootp flags: 0x000 (Unicast)	
Client IP address: 0.0.0.0 (0.0.0.0)	(64 OCTETS)
Your (client) IP address: 0.0.0.0 (0.0.0.0)	28 OCTETS)
Next server IP address: 0.0.0.0 (0.0.0.0)	20 001213)
Relay agent IP address: 0.0.0.0 (0.0.0.0)	IABLE)
Client MAC address: Grandstr_01:fc:42 (00:0b:82:01:fc:42	

N	0	Len	Time	Source	Destination	Protocol	Info	
	1	314	0.000000	0.0.0.0	255.255.255.255	DHCP	DHCP Discover	- Transaction ID
	2	342	0.000295	192.168.0.1	192.168.0.10	DHCP	DHCP Offer	- Transaction ID
	3	314	0.070031	0.0.0.0	255.255.255.255	DHCP	DHCP Request	- Transaction ID
	- 4	342	0.070345	192.168.0.1	192.168.0.10	DHCP	DHCP ACK	- Transaction ID

Request and ACK

	DHCPF	REQUEST		_		D	HCPACK		
UDP Src=0.0.0.0 sPort=68 Dest=255.255.255.255 dPort=67				Client uses Port 68	UDP Src=192.168.1.1 sPort=67 Dest=255.255.255.255 dPort=68				
OP	HTYPE	HLEN	HOPS		OP	HTYPE	HLEN	HOPS	
0x01	0x01	0x06	0x00		0x02	0x01	0x06	0x00	
	3	KID					XID		
0x3903	F326				0x3903F3	326			
	SECS	FL	AGS			SECS		FLAGS	
0x0000		0x0000			0x0000	0x0000 0:		0x0000	
(CIADDR (Cli	ent IP Add	ress)	Assuming the client is choosing the	CIADDR (Client IP Address)				
0x0000	0000				0x0000000				
	YIADDR (Yo	ur IP Addr	ess)		YIADDR (Your IP Address)				
0x0000	0000			address form	0xC0A80	0xC0A80164 Client's NEW IP Address			
	SIADDR (Ser	ver IP Add	ress)	the server	SIADDR (Server IP Address)				
0xC0A	80101			Server IP Address	0xC0A80101				
G	IADDR (Gate	way IP Ad	dress)		GIAD	GIADDR (Gateway IP Address switched by relay)			
0x0000000			_	0×0000000					
СНА	DDR (Client	Hardware	Address)	-		CHADDR (Clie	ent Hardware A	Address)	
0x0005	3C04			-	0x000530		vor		
0-0050		lient		-	0x8D590	000	VEL		

When the DHCP server receives the DHCPREQUEST message from the client, the configuration process enters its final phase. The acknowledgement phase involves sending a DHCPACK packet to the client.



- if a host does not know its IP address, the server uses field YOUR IP ADDRESS to supply the value
- server uses fields SERVER IP ADDRESS and SERVER HOST NAME to give the host information about the location of a server
- ROUTER IP ADDRESS contains the IP address of a default router
- DHCP allows a computer to negotiate to find a **boot image**
 - The computer is boot up, request and OS
 - the host fills in field BOOT FILE NAME with a request
 - The DHCP server does not send an image
 - The host uses TFTP

Early Release

- The user can end the lease through a process called early lease termination or lease release
- This is a very simple, unidirectional communication
 - The client sends a special DHCPRELEASE message unicast to the server that holds its current lease
 - The server then records the lease as having been ended
 - It does not need to reply back to the client (no ACK)
- Client can just assume that the lease termination has been successful
- Having clients send DHCPRELEASE to end a lease is considered a *courtesy*, rather than a requirement
- DHCP servers are designed to handle the case where a client "disappears" without formally ending an existing lease
 - Sending a DHCPRELEASE is clearly more efficient, however!

Indirect DHCP Server Access



- DHCP broadcasts on the local network to find a server
- DHCP does not require each individual network to have a server
 - Instead, a DHCP relay agent forwards requests and responses between a client and the server
- At least one relay agent must be present on each network
 - The relay agent must be configured with the address of the appropriate DHCP server
- When the DHCP server responds
 - The relay agent forwards the response to the client

Network Address Translation (NAT)

Addresses are growing! What is the solution?

- Use IPV6
- Use NAT
- □ NAT:
 - Allows using one IP address per company
 - Internally new addresses can be added!

How?

NAT Operation

- □ IP reserved addresses
 - 10.x.y.z
 - 172.16.x.y
 - 192.168.x.y
- Receiving a packet from the Internet
 - Sender
 - Add IP address
 - TCP will have the destination port (0-1023) – standard port group
 - The port determines which server on the remote (destination) side to process the packet
 - NAT box:
 - Using the PORT address in TCP, change the IP address to a designated address (10.0.01)
- Sending a packet into the Internet
 - NAT box:
 - Changes the source address from 10.0.0.1 to 198.60.42.12



NAT Issues...

- □ Addresses are not unique: many 10.0.0.1!
- NAT controls the incoming and outgoing packets – reliability!
- NAT accesses TCP and IP layers layers should work independent of one another
- NAT only allows TCP/IP or UDP/IP
- NAT does not support applications which insert the IP address in the body (FTP or H.323)

DEMO http://www3.rad.com/networks/2005/prvt-nat/main2.htm



Remember...

- This is My MAC; what is my IP address? RARP / DHCP
- This is the destination host name, what it is IP address? DNS Server
- This is the IP address, what is your MAC address? ARP

References

Tanenbaum
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