Internet Protocols

Addressing & Services

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Virtual vs. Physical Networks

- MAC is the part of the underlying network
 - MAC is used on the LAN
- What is the addressing mechanism in WAN?
 - WAN is interconnections of man many LANS
- Networking addressing is required
 - Making the network of networks to appear seamless
- For Internet we use IP addressing

TCP/IP Stack		
	ТСР	
	IP	
	LLC	
	MAC	
	Physical	

Ethernet Addressing

□ MAC address is 48 bits:

- 24 bits (OUI Organizationally unique Identifier
- 24 bit hardware address burned in the ROM



Ethernet Addressing

These MAC addresses are found via:	
http://standards.ieee.org/regauth/oui/index.shtml	

Enter MAC:

Submit Query



Network Layer Architecture

- Layer 3 of the seven-layer
- Provides services to upper layer (Primitives and parameter)
- The Network Layer is responsible for routing packets delivery
 - Note the Data Link Layer is responsible for Media Access Control, Flow Control and Error Checking
- Connection model: connectionless communication
 - No setup path is required
 - The recipient does not have to send an acknowledgement
- Provides unique host addressing

CP/IP Stack		
	ТСР	
	IP	
	LLC	
	MAC	
	Physical	

Network Layer Examples

- IPv4/IPv6, Internet Protocol
- DVMRP, Distance Vector Multicast Routing Protocol
- ICMP, Internet Control Message Protocol
- IGMP, Internet Group Multicast Protocol
- PIM-SM, Protocol Independent Multicast Sparse Mode
- PIM-DM, Protocol Independent Multicast Dense Mode
- IPsec, Internet Protocol Security
- IPX, Internetwork Packet Exchange
- RIP, Routing Information Protocol
- DDP, Datagram Delivery Protocol
- BGP, Border Gateway Protocol



Internet Protocol

- We focus on IP
- IP was the connectionless datagram service
- Originally introduced by Vint Cerf and Bob Kahn in 1974 to be interfaced with TCP
 - The first major version of IP is known as Internet Protocol Version 4 (IPv4) – dominant
 - Internet Protocol Version 6 (IPv6) is the successive version
- Main responsibility: addressing hosts and routing datagrams (packets) from a source host to the destination host across one or more IP networks
 - Addresses identify hosts
 - Provides a logical location service

Internet Protocol

- IP Design Issues
 - Interconnection
 - Routing
 - Static or Dynamic
 - Looping and lifetime
 - Fragmentation
 - Error Control
 - Flow Control
 - IP Header and Addressing

IP Encapsulation in Frames

- The IP datagram contains data and IP address
- The IP datagram is encapsulated in a frame with physical address
- The header changes as the frame goes from one network domain to the next

IP Datag	ram		
Frame			
	Physical Address	IP Address	Data



IP: Connectionless Internetworking

Advantages

- Flexibility and robust
- No unnecessary overhead

Unreliable

- Not guaranteed delivery (no ACK is required)
- Not guaranteed order of delivery

Packets can take different routes

 Reliability is responsibility of next layer up (e.g. TCP)

IP Routing

End systems and routers maintain routing tables

- Indicate next router to which datagram should be sent
- Static
 - May contain alternative routes
- Dynamic
 - Flexible response to congestion and errors
- Source routing
 - Source specifies route as sequential list of routers to be followed

IP Datagram Lifetime

- Datagrams could loop indefinitely
 - Consumes resources
 - Transport protocol may need upper bound on datagram life
- Datagram marked with lifetime
 - Time To Live field in IP
 - Once lifetime expires, datagram discarded (not forwarded)
 - Hop count
 - Decrement time to live on passing through a each router
 - Time count
 - Need to know how long since last router

IP Packet TTL

- TTL (time-to-live) refers to the number of router hops the IP packet is allowed before it must be discarded.
 - Each router that receives a packet subtracts one from the count in the TTL field.
 - When the count reaches zero, the router detecting it discards the packet and sends an Internet Control Message Protocol (ICMP) message back to the originating host



IP Flow Control and Error Control

Error Control

- FCS (frame Check Sequence)
- CRC
- Congestion errors / lifetime expiration
- Error notification may not possible error in address!

Flow Control

- ICMP (change the sending rate)
- Node-to-node backoff

		- · · ·
IP	TCP	
Header	Header	TCD novload (394 actets)
(20	(20	1 CP paytoat (384 octets)
octets)	octets)	

IP Services and Versions

- Part of TCP/IP
 - Used by the Internet
- Specifies interface with higher layer
 - e.g. TCP
- Specifies network protocol format and mechanisms
- IPv4
 - Addresses are 32 bits wide
 - Its header is 20 bytes at minimum
 - Uses doted-decimal notation (e.g. 43.23.43.56) using octets
- □ IPv6
 - Provides larges address domain; addresses are 128 bits wide
 - Multiple separate headers are supported
 - Offers roaming features
 - Handles audio and video; providing high quality paths
 - Supports unicast, multicast, anycast

IPv4 Addressing

Internet Addressing

- Over half million networks are connected to the Internet 5 billion users by 2015!
- Network numbers are managed by ICANN (Internet Corporation for Assigned Names and Numbers) -<u>http://www.icann.org/</u>
 - Delegates part of address assignments to regional authorities called registrars
 - Registrars are authorized by ICANN to assign blocks of addresses
 - IP address blocks are given to ISPs and companies
 - ISPs distribute individual addresses to users and organizations
- IP addresses are based on dotted decimal notation: 192.41.7.32 (Octets from 0 to 255 – 8 bits)
 - IP address 0.0.0.0 refer to machine's own network when it is being booted (This host)
 - 255.255.255.255 broadcast on the LAN
 - 127.x.y.z reserved for loopback testing

IP Addressing – IPv4

- A network IP address is divided into Netid and Hostid
- Also called Prefix and Suffix .
- IP Address classification

Class	Leading bits	Network Address (Netid)	Host Address (Hostid)
Class A	0	7 bit (125)*	24 bit (16,777,216)
Class B	10	14 bit (16,382)	16 bit (65,534)
Class C	110	21 bit (2,096,150)	8 bit (254)
Class D (multicast)	1110	Multi	cast Address
Class E (reserved)	1111	Reserved (not assigned)	

- Some values are reserved (e.g., all zero, all one)!
- Leading bits refer to most significant bits

IP Addressing Classification





Example of IP Addressing

Q1: Determine the network address for the following IP addresses:

```
1- 84.42.58.11 (84 = 54 Hex = 0101 \ 0100)

→ Netid=84.0.0.0
```

- \rightarrow Class A
- → Hostid=0.42.58.11
- 2- 144.54.67.5 (144 = 90 Hex = 1001 0000)
 - → Netid=144.54.0.0
 - \rightarrow Class B
 - → Hostid=0.0.67.5

Q2: What type of IP address classification will a large organization with 1000 individual users in 150 dispersed buildings use? \rightarrow Class B



Reserved Addresses

Some addresses are reserved

Loopback 127.0.0.1 commonly used for Loopback

- When Loopback address is used packet do not penetrate to the network (used to check the network card)
- If HostID is all-one → packets are broadcasted to all the hosts on the network
 - Hardware must support broadcast delivery otherwise software must send single messages to each host
- In case of BSC (BSD Berkeley Software Distribution <u>http://www.bsd.org/</u> when HostID is all-zero → packets are broadcasted to all hosts on the network
 - BSD one of the original Unix Distributions
 - Implemented TCP/IP
 - Many are still using it

Subnets and Subnet Masks

- Each LAN assigned subnet number
- Host portion of address partitioned into subnet number and host number
- Local routers route within the subnet
- Subnet mask indicates which bits are subnet number and which are host number
 - Ones indicate NetID
 - Zeros indicate Hosts
- Insulate overall internet from growth of network numbers and routing complexity

R1

Net ID/Subnet ID: 192.228.17.32 00 10 00 00 Subnet number: 1



Network A



IP Address: 192.228.17.33 Host number: 1 IP Address: 192.228.17.57 Host number: 25 25+32=57

Router LAN1 (subnet1)



Routing Using Subnets



Host number: 1



Subnet Mask Example

Given IP Address of LAN X: 192.228.17.32 and its MASK is defined as 255.255.255.224, will 192.228.17.58 be a valid address on the LAN?

If so, what will be its Host Number?

Classes and Subnets...

- Classful routing is not very efficient
- Having IP address classes creates issues
 - Addresses can be under utilized (Class A)
 - Addresses can be over utilized (Class C)
 - Management of addresses may be difficult
 - Organizations can grow!
- **Example:** Site 1: 12 Hosts / Site 29 Hosts
 - We use 128.211.0.0 → C Class; Dedicated 256-2=254 users! But only few are used
- Alternatives
 - Subnets addressing
 - Classless addressing
- Classless Inter Domain Routing (CIDR)
 - Allocate remaining IP addresses in variable-sized block; no regard to class!
 - Use 32-bit mask!
 - Uses a single routing table

Classless vs. Classfull



More about subnets....

- Routers can be connected to multiple LANs
- LANS are divided into subnets each identified by a subnet mask: 255.255.252.0 (... 1111 1100 0000 0000) → 32-bit-10-bit=22-bit to identify the subnet!
 - Mask: netID + SubnetID or /22 (subnet mask is 22 bit long) we mask or "hide" the first 22 bits
 - Subnets are <u>not visible</u> outside the network

Represents the NetId part in the mask

- Example: Assume subnet mask is 255.255.252.0/22⁴
 - Subnet 1: 130.50.4.0
 - Subnet 2: 130.50.8.0→000010|00 0000000
 - Subnet 3: 130.50.12.0→

...000010|00 0000000 ...000011|00 00000000

- Assume a packet's destination is 130.50.15.6 → which subnet does it belong to?
- Mask: ... 1111 11|00 0000 0000
- Adrs: ... 0000 11|11 0000 0110
- 0000 11|00 0000 0000
- Hence: the packet must go to Subnet 3 (130.50.12.xx)



		Length (CIDR)	Address Mask	Notes
		/0	0.0.0.0	All De (equivalent to no mask)
		/1	128.0.0.0	
		/2	192.0.0.0	
	CIDR Notation	/3	224 . 0 . 0 . 0	
		/4	240 . 0 . 0 . 0	
_		/5	248 . 0 . 0 . 0	
		/6	252 . 0 . 0 . 0	
	Example 1:	/7	254 . 0 . 0 . 0	
	Calculate the mask for	/8	255 . 0 . 0 . 0	Original Class A mask
	102 168 100 0/24	/9	255 . 128 . 0 . 0	
	192.100.100.0/24	/10	255 . 192 . 0 . 0	
		/11	255 . 224 . 0 . 0	
	255.255.255.0	/12	255 . 240 . 0 . 0	
		/13		
		/ 14		
	Example 2:	/15		
	Assuming a host (connection)	/17		Original Class B mask
	has an address of 172 16 45 0	/18	255 255 102 0	
	With most value of 2EE 2EE 0	/19		
		/20	255 255 240 0	
	What will be the network address?	/21	255 255 248 0	
		/22	255 255 252 0	
	9 hits are for HostID	/23	255 . 255 . 254 . 0	
		/24	255 . 255 . 255 . 0	Original Class C mask
	45→ <u>00 10 11 0</u> 1	/25	255 . 255 . 255 . 128	enginal olabe o mask
	$00\ 10\ 11\ 0 \rightarrow 22 \rightarrow \text{NetId}:\ 192.16.22.0$	/26	255 . 255 . 255 . 192	
	$172 \rightarrow 1010 \ 1100 \rightarrow B$	/27	255 . 255 . 255 . 224	
		/28	255 . 255 . 255 . 240	
		/29	255 . 255 . 255 . 248	
		/30	255 . 255 . 255 . 252	
		/31	255 . 255 . 255 . 254	
		/32	255 . 255 . 255 . 255	All 1s (host specific mask)

Reserved Addresses

- **10.0.0.**/8
- **169.254.0.0/16**
- **172.16.0.0/12**
- **192.168.0.0/16**

Classless Routing - Example

	<u> </u>			·
University	First address	Last address	How many	Written as
Cambridge	194.24.0.0	194.24.7.255	2048	194.24.0.0/21
Edinburgh	194.24.8.0	194.24.11.255	1024	194.24.8.0/22
(Available)	194.24.12.0	194.24.15.255	1024	194.24.12/22
Oxford	194.24.16.0	194.24.31.255	4096	194.24.16.0/20

Address	Mask
C: 11000010 00011000 0000000 0000000	11111111 11111111 11111000 00000000
E: 11000010 00011000 00001000 00000000	11111111 11111111 11111100 00000000
O: 11000010 00011000 00010000 00000000	11111111 1111111 11110000 00000000

What happens if a packet has an address of 194.24.17.4? Where does it go?

- O:1111 0000 0000 0000 AND
 -0001 0001 0000 0100 \rightarrow

→0001 0000 0000 0000 (194.24.16.0)→ <u>Oxford</u>

How do you represent class B using CIDR?

16-bit NetID + 16-bit HostID \rightarrow /16

2048=2^11

8bits+3bits

32-11=21

Routing Decisions by the Router

□ If (Mask [i] & D) == Destination [i] \rightarrow Forward next hop



Each port has a diff. address

Destination	Mask	Next Hop
30.0.0.0	255.0.0.0	40.0.0.7
40.0.0.0	255.0.0.0	deliver direct
128.1.0.0	255.255.0.0	deliver direct
192.4.10.0	255.255.255.0	128.1.0.9

Assume R2 receives packet with destination 192.4.10.3!

Router IP Address / Connections

- IP addresses refer to connections
- The suffixes for each router can be the same for ease of remembering



Note:

Route 2 uses the same suffix (suffix in this case is the last byte) Router 1 Uses different suffix

Questions:

-Can you reach 78.0.1.17?

-Router 1 passes a packet with destination address

131.108.255.255 ;where does it go?

-If BSD is used what will be the broadcast address 131.108.0.0



Practice Problems:

Provide the following parameter values for each of the network classes A, B, and C. Be sure to consider any special or reserved addresses in your calculations.

- I. Number of bits in network portion of address
- b. Number of bits in host portion of address
- Number of distinct networks allowed
- d. Number of distinct hosts per network allowed
- •. Integer range of first octet

→ What percentage of the total IP address space does each of the network classes represent?

What is the difference between the subnet mask for a Class A address with 16 bits for the subnet ID and a class B address with 8 bits for the subnet ID?

Is the subnet mask 255.255.0.255 valid for a Class A address?

- \rightarrow Given a network address of 192.168.100.0 and a subnet mask of 255.255.255.192,
 - a. How many subnets are created?
 - **b.** How many hosts are there per subnet?

References

Tanenbaum
Tomasi Text Book
Comer Text book